

Princess Nora University
Faculty of Computer & Information Systems



جامعة الأميرة نورة بنت عبد الرحمن
Princess Nora Bint Abdul Rahman University

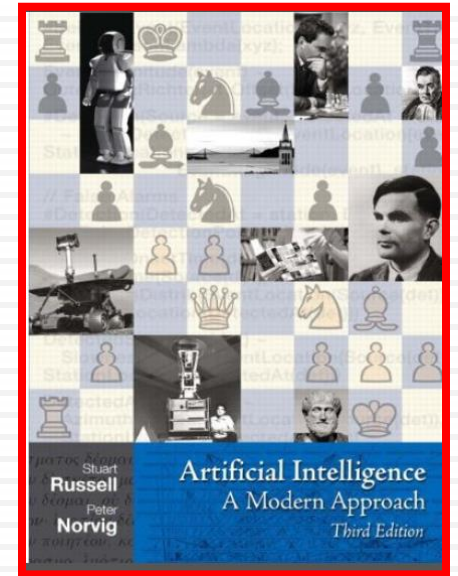


ARTIFICIAL INTELLIGENCE

(CS 370D)



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(CHAPTER-6)

CONSTRAINT SATISFACTION PROBLEMS



Outline

- What is a CSP
- CSP applications
- Backtracking search
- Problem structure and decomposition





Mean of Constraint Satisfaction Problems





Constraint Satisfaction Problems

- A constraint satisfaction problem consists of three components,

X , D , and C

X is a set of variables, $\{X_1, \dots, X_n\}$.

D is a set of domains, $\{D_1, \dots, D_n\}$, one for each variable.

C is a set of constraints that specify allowable values.

- Each constraint C_i limits the values that variables can take,

- e.g., $V_1 \neq V_2$



A *state* is defined as an *assignment* of values to some or all variables.





CSPs (continued)

- An assignment is **complete** when every variable is mentioned.
- A **solution** to a CSP is a **complete** assignment that satisfies all constraints.
- Some CSPs require a solution that maximizes an *objective function*.





Varieties of constraints

- **Unary constraints** involve a single variable.
 - e.g. $SA \neq green$

- **Binary constraints** involve pairs of variables.
 - e.g. $SA \neq WA$

- **Higher-order constraints** involve 3 or more variables.
 - Professors A, B, and C cannot be on a committee together
 - Can always be represented by multiple binary constraints

- **Preference** (soft constraints)
 - e.g. *red* is better than *green* often can be represented by a cost for each variable assignment
 - combination of optimization with CSPs





Constraint Satisfaction Applications





CSPs (continued) Examples of Applications

□ Scheduling :

- Scheduling the timetable of lectures
- Airline schedules
- Scheduling your MS or PhD thesis exam 😊

□ Configuration

- To arrange objects of different sizes in a container, Map Coloring.

- CSP solvers can be faster than state-space searchers because the CSP solver can quickly eliminate large swatches of the search space.





Example problem – timetable of university (1)

- state - collection of schedule objects
 - ▣ schedule
 - professor ($\sim 10^2$) (~ 350 at LU)
 - course ($\sim 10^3$)
 - time slot (~ 10) (16 at LU)
 - location ($\sim 10^2$)
- constraints
- objective function





Example problem – timetable of university (1)

□ Constraints

- professors assigned to courses
- one course per location per time slot
- location capacity \geq course (projected) enrollment
- professors assigned only once per time slot
- rooms assigned only once per time slot
- courses in program in different time slots





Example problem – timetable of university (1)

□ objective function

- minimize total course enrollment in early morning time slots
- distribute program courses across days
- accommodate special requests





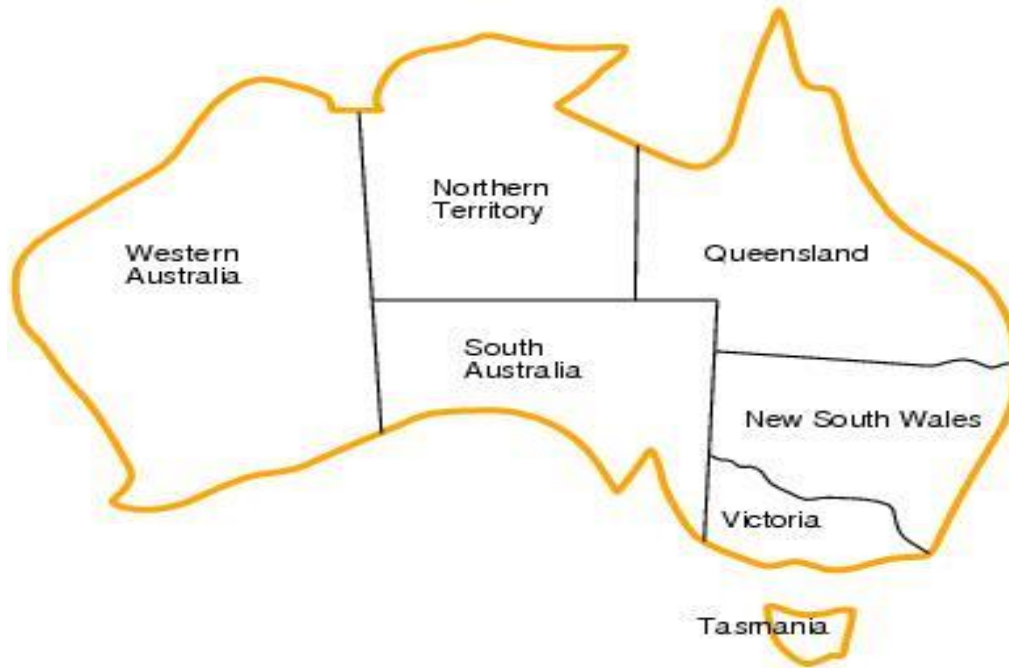
Solving CSPs

- analysis of problem state structure is critical
- **two basic approaches:**
 - ▣ start state is null: assign state variables one at a time and backtrack if constraint is violated
 - ▣ assign variables 'randomly/intelligently' to make a start state and use greedy search by changing variable assignments





CSP example: map coloring

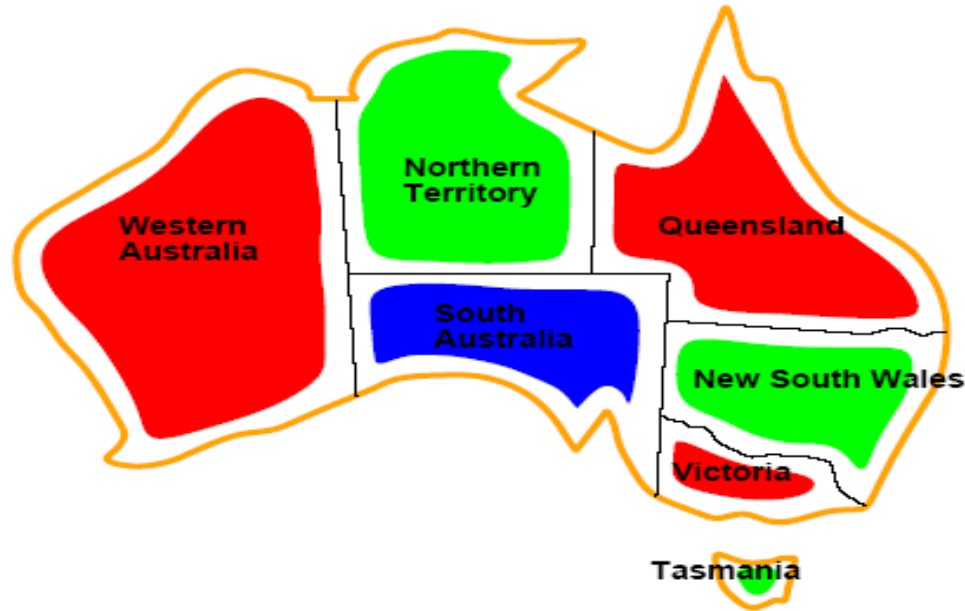


- Variables: WA, NT, Q, NSW, V, SA, T
- Domains: $D_i = \{red, green, blue\}$
- Constraints: adjacent regions must have different colors.
 - E.g. $WA \neq NT$





CSP example: map coloring



□ Solutions are assignments satisfying all constraints, e.g.

$\{WA=red, NT=green, Q=red, NSW=green, V=red, SA=blue, T=green\}$





Solving CSPs

- As general searching problem.
- Back-Tracking.
- Arc-consistency Checking.
- Heuristics.





As General Searching Problem





CSP as a standard search problem

- A CSP can easily be expressed as a standard search problem.

- **Incremental formulation**
 - **Initial State**: the empty assignment $\{\}$

 - **Successor function**: Assign a value to any unassigned variable provided that it does not violate a constraint

 - **Goal test**: the current assignment is complete
(by construction its consistent)

 - **Path cost**: constant cost for every step (not really relevant)





Backtracking Search





Backtracking Search

- Variable assignments are **commutative** if the order of any given set of actions has no effect on the outcome.
 - i.e., [WA = red then NT = green] same as [NT = green then WA = red].
- Only need to consider assignments to a single variable at each node. and backtracks when a variable has no legal values left to assign.
- Depth-first search for CSPs with single-variable assignments is called **Backtracking Search**:



- ✓ Chooses values for one variable at a time.
- ✓ Checks for constraint violations at each assignment.



- ✓ Backtracks when a variable has no legal values left to assign.



Backtracking search

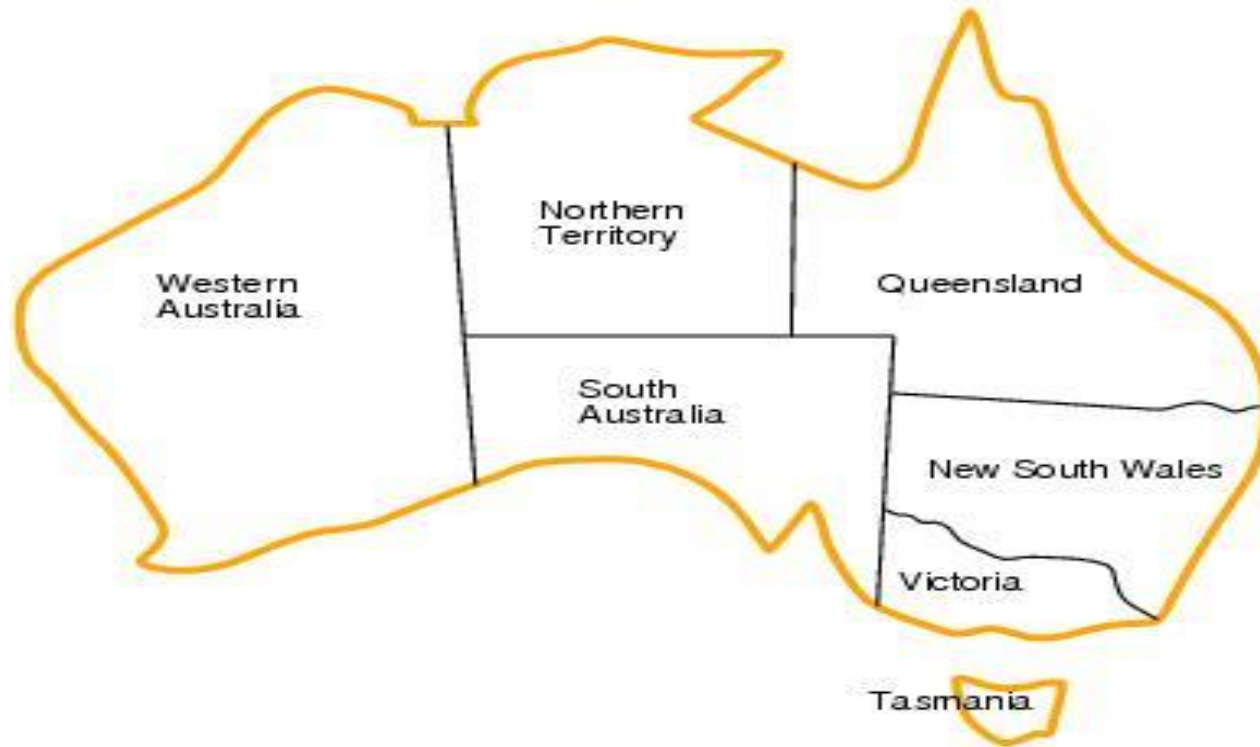
```
function BACKTRACKING-SEARCH(csp) return a solution or failure  
  return RECURSIVE-BACKTRACKING({}, csp)
```

```
function RECURSIVE-BACKTRACKING(assignment, csp) return a solution or failure  
  if assignment is complete then return assignment  
  var ← SELECT-UNASSIGNED-VARIABLE(VARIABLES[csp],assignment,csp)  
  for each value in ORDER-DOMAIN-VALUES(var, assignment, csp) do  
    if value is consistent with assignment according to CONSTRAINTS [csp]  
    then  
      add {var=value} to assignment  
      result ← RRECURSIVE-BACKTRACKING(assignment, csp)  
      if result ≠ failure then return result  
      remove {var=value} from assignment  
  return failure
```



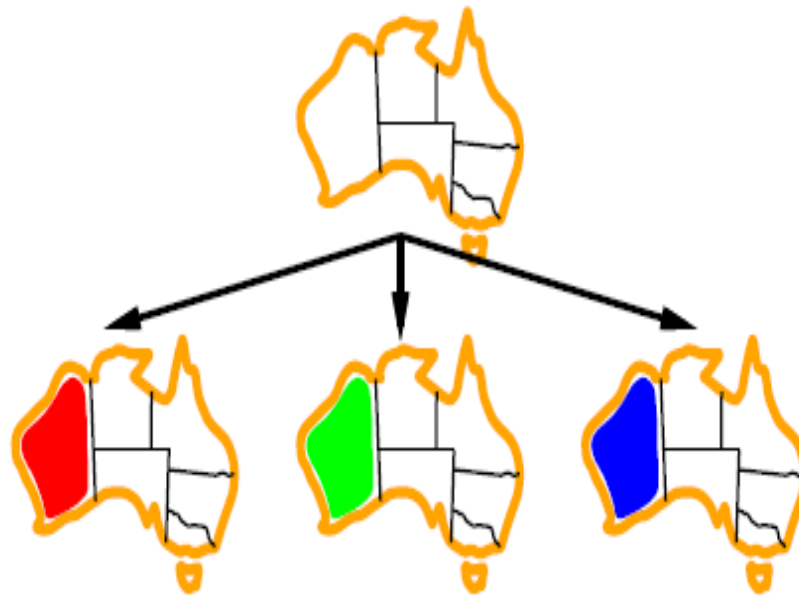


Backtracking example

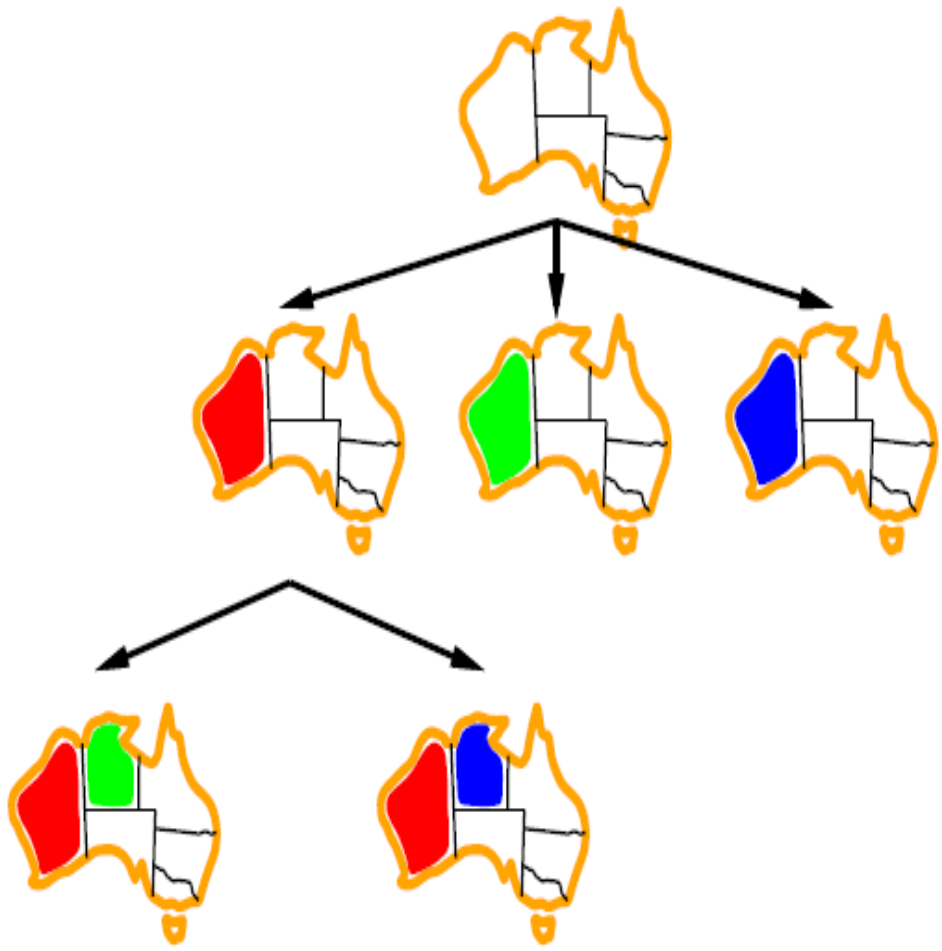




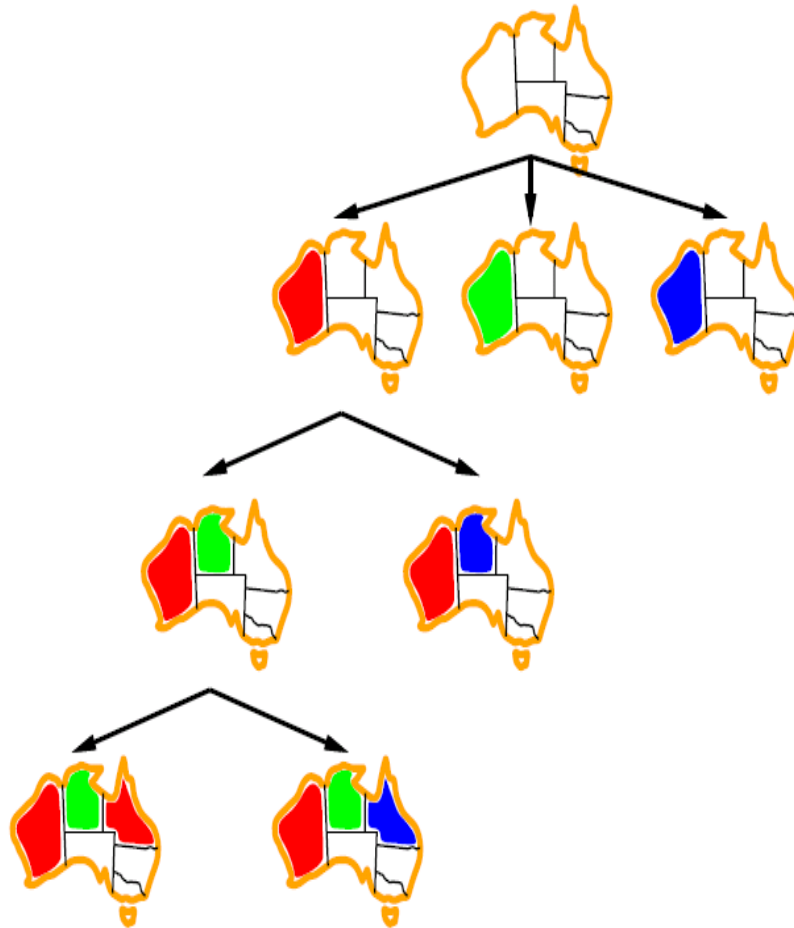
Backtracking example



Backtracking example



Backtracking example



Constraints on SA will eventually cause failure when $WA \neq Q$. When not the same color (bottom right), SA cannot be assigned.

The algorithm will backtrack to a node with unexplored states.

For example, **such as WA=red, NT=blue.**





Improving Backtracking Efficiency

□ Consider

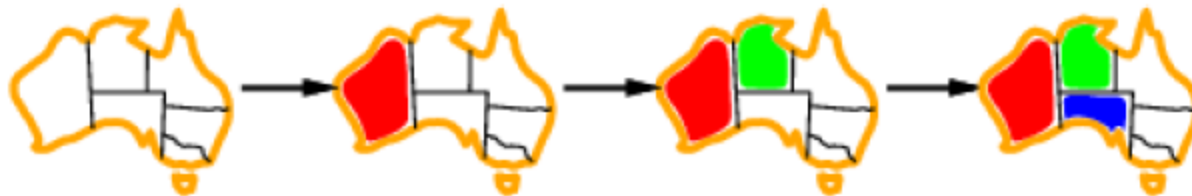
- Which variable should be assigned next?
- In what order should the variables be tried?
- Can we detect inevitable failure early?
- Can we take advantage of problem structure?





Most Constrained Variable or minimum remaining values (MRV)

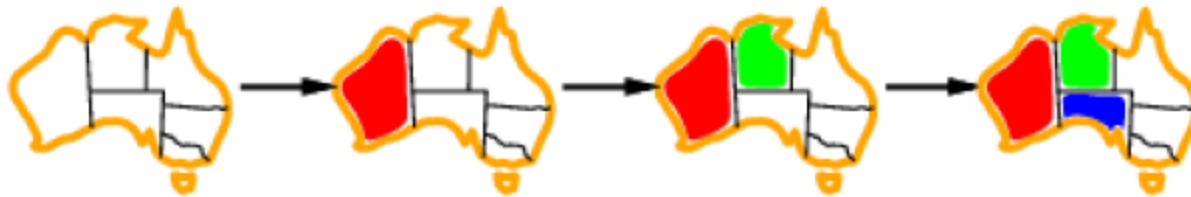
- **Heuristic Rule:** Choose the variable with the **fewest** legal values.
- **Idea** is that variable with *fewest* choices will likely fail sooner allowing branch to be *pruned* early.
 - Legal values are those that do not conflict with any already assigned to a neighbor (i.e. do not violate constraints).
 - After WA=red, NT and SA have only 2 remaining legal values, the other unassigned variables all have 3.
 - **Choose between NT and SA.**





□ Most constraining variable

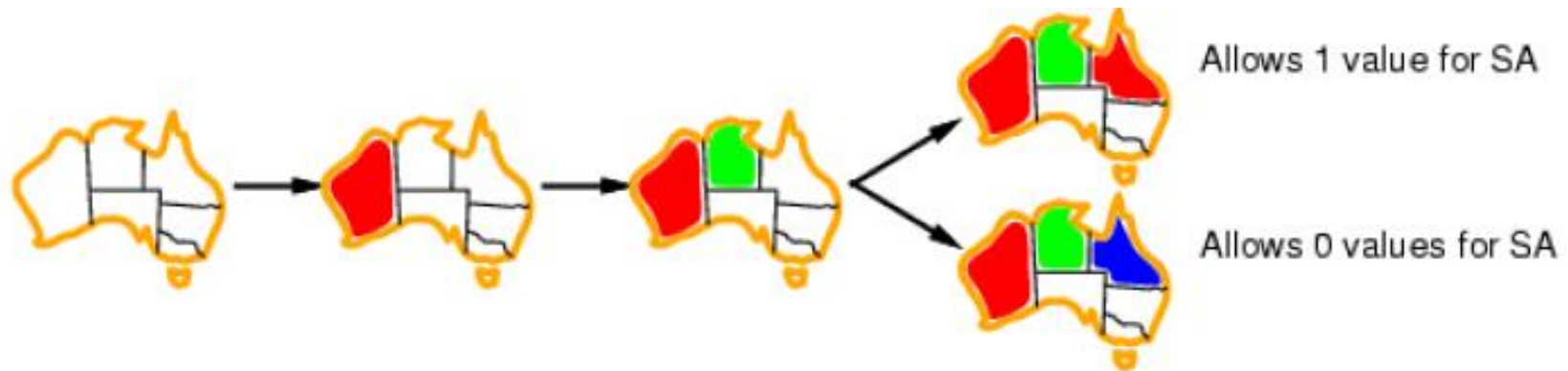
- Choose the variable that constrains the most remaining variables
- SA places constraints on 5 variables; WA, NT, Q, V, NSW only on 3, T on 0.
- Purpose is to prune early by reducing number of choices.
- Tie breaker among minimum remaining variables (MRV)





Least Constraining Value

- **Heuristic Rule:**
- Given a variable, choose the least constraining value



Idea is to explore paths, those with most options, most likely to lead to **a solution**.





Forward checking



WA

NT

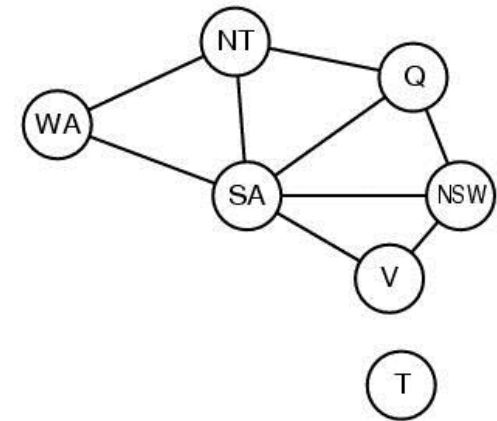
Q

NSW

V

SA

T

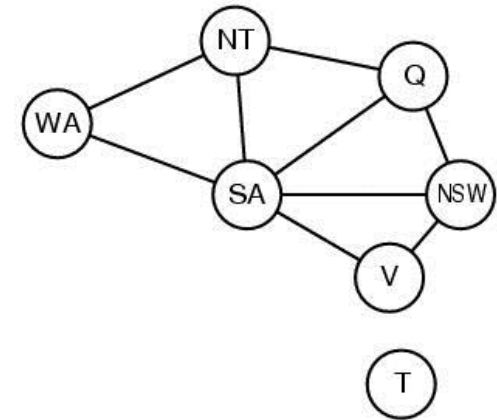


- Can we detect inevitable failure early?
 - ▣ *And avoid it later?*
- **Forward checking idea:** keep track of remaining legal values for unassigned variables.
- Terminate search when any variable has **no legal values.**





Forward checking



- Assign $\{WA=red\}$
- Effects on other variables connected by constraints to WA



- *NT can no longer be red*
- *SA can no longer be red*

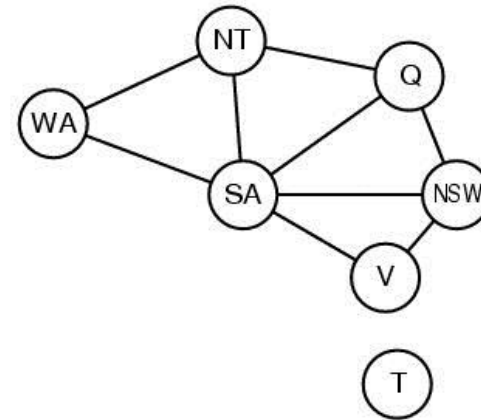




Forward checking



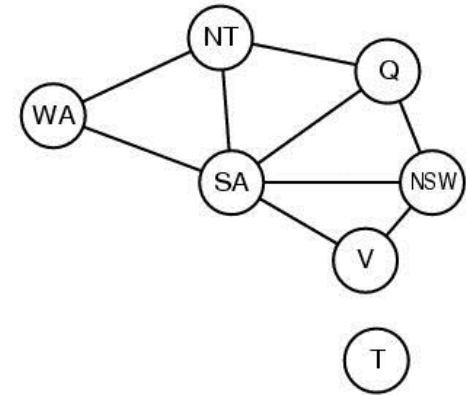
WA	NT	Q	NSW	V	SA	T
Red	Green	Blue	Red	Green	Blue	Red
Red	Green	Blue	Red	Green	Blue	Red
Red	Green	Blue	Red	Green	Blue	Red



- Assign {Q=green}
- Effects on other variables connected by constraints with WA
 - NT can no longer be green
 - NSW can no longer be green
 - SA can no longer be green



Forward checking



WA	NT	Q	NSW	V	SA	T
Red Green Blue	Red Green Blue	Red Green Blue	Red Green Blue	Red Green Blue	Red Green Blue	Red Green Blue
Red	Green Blue	Red Green Blue	Red Green Blue	Red Green Blue	Green Blue	Red Green Blue
Red	Blue	Green	Red Blue	Red Green Blue	Blue	Red Green Blue
Red	Blue	Green	Red	Blue		Red Green Blue

- If V is assigned *blue*
- Effects on other variables connected by constraints with WA
 - ▣ NSW can no longer be blue
 - ▣ SA is *empty*



□ FC has detected that partial assignment is *inconsistent* with the constraints and backtracking can occur.

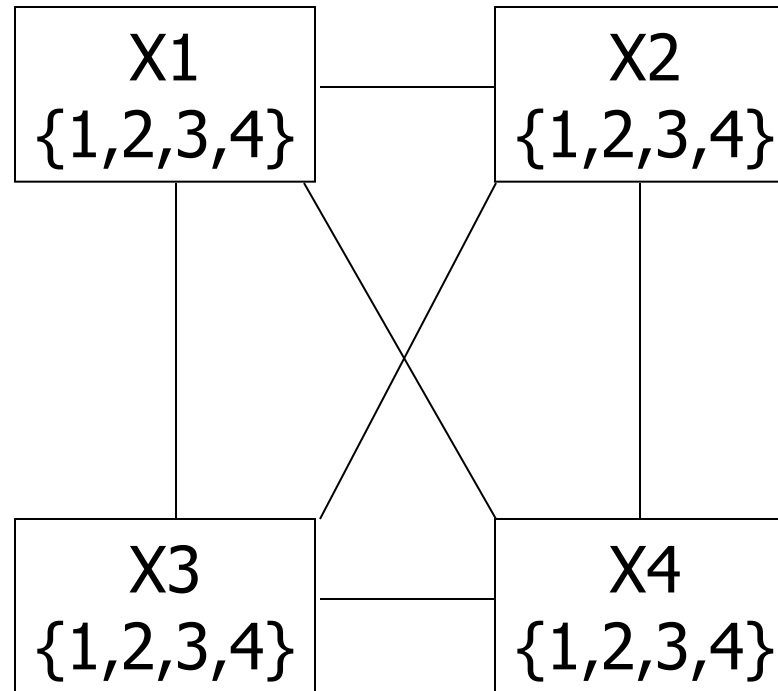




Forward checking

Example: 4-Queens Problem

	1	2	3	4
1		■		■
2	■		■	
3		■		■
4	■		■	

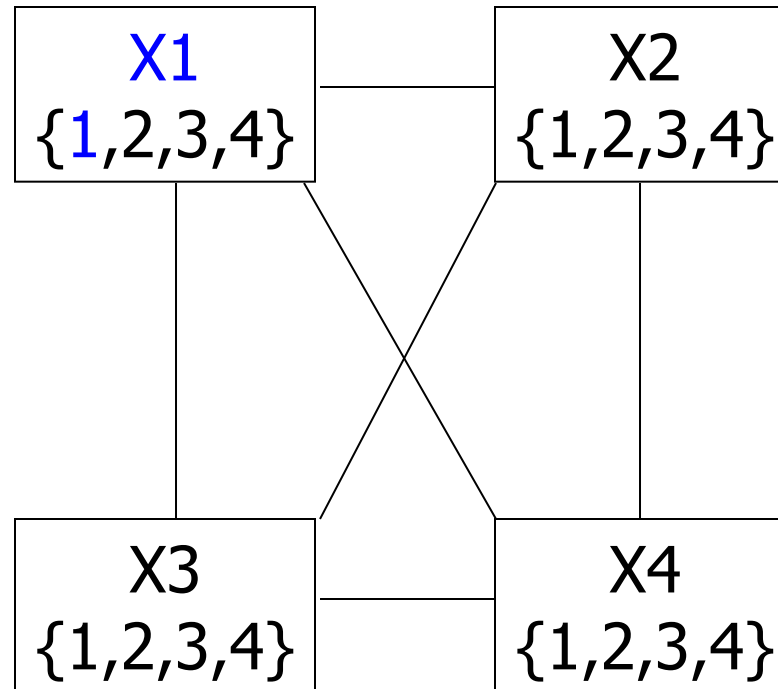




Forward checking

Example: 4-Queens Problem

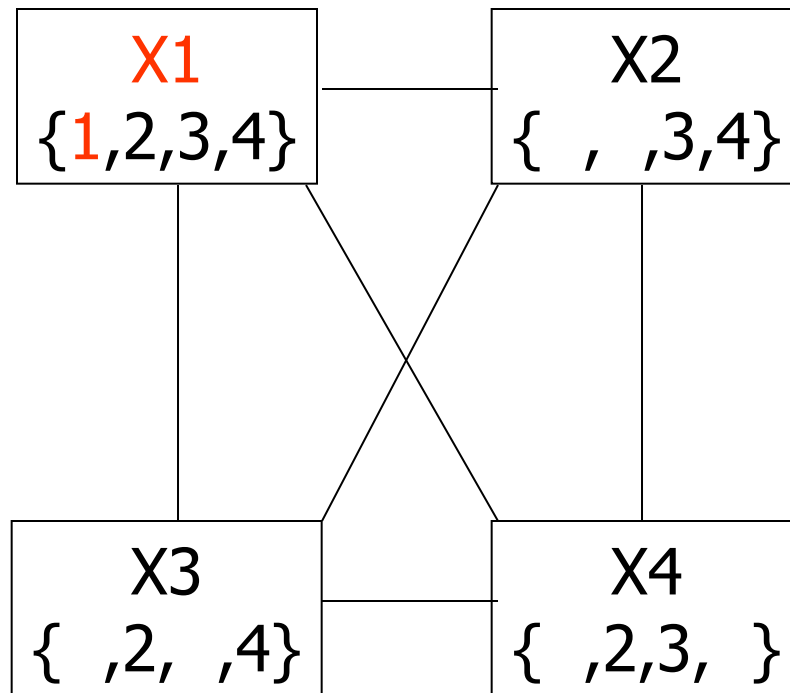
	1	2	3	4
1		■	□	■
2	■	□	■	□
3	□	■	□	■
4	■	□	■	□





Example: 4-Queens Problem

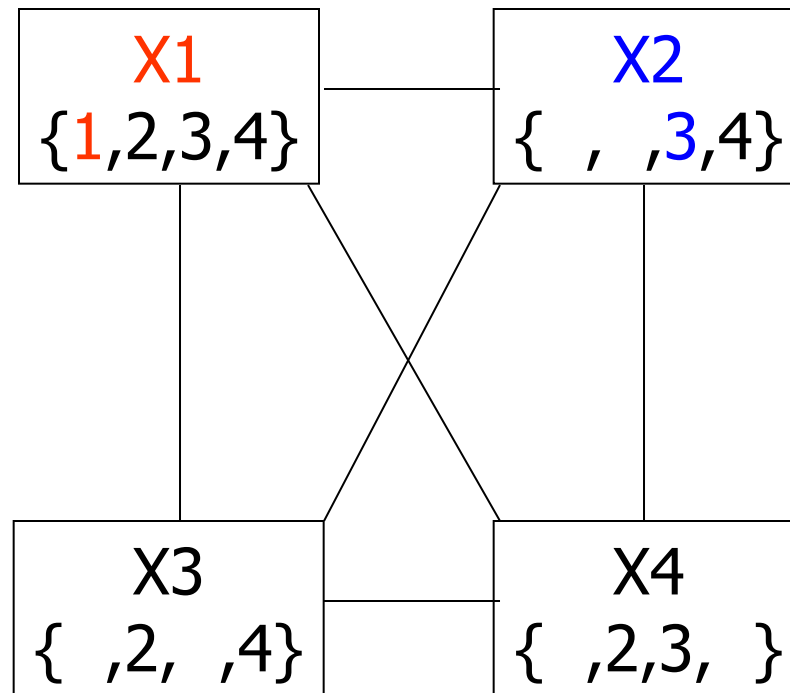
	1	2	3	4
1				
2				
3				
4				





Example: 4-Queens Problem

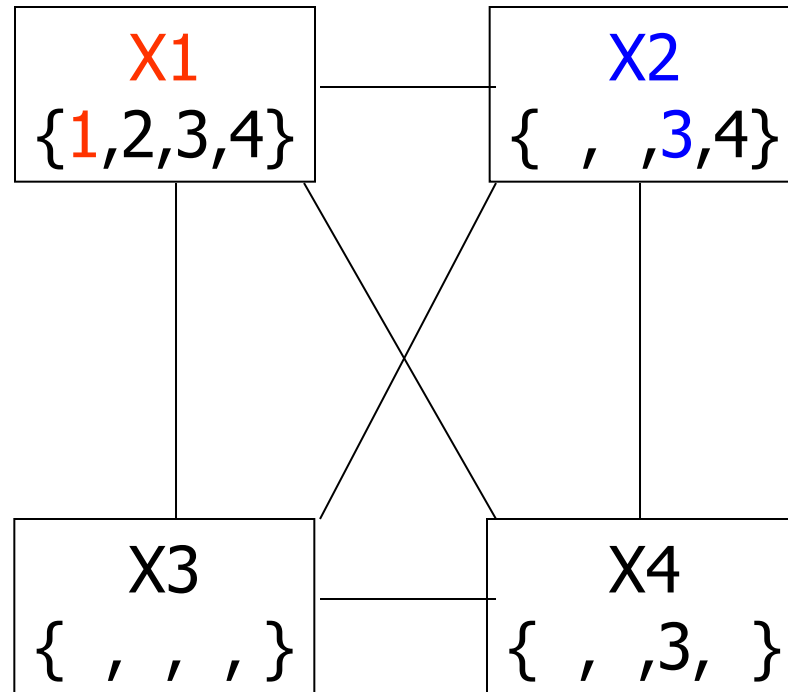
	1	2	3	4
1	★	●	●	●
2		●		
3		★	●	
4				●





Example: 4-Queens Problem

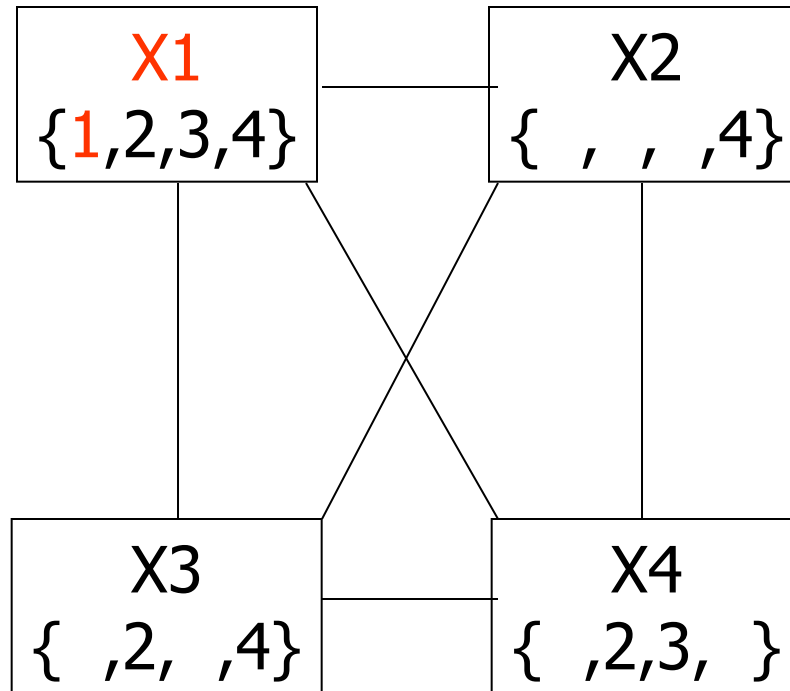
	1	2	3	4
1	★	●	●	●
2	■	●	●	□
3	□	★	●	●
4	■	□	●	●





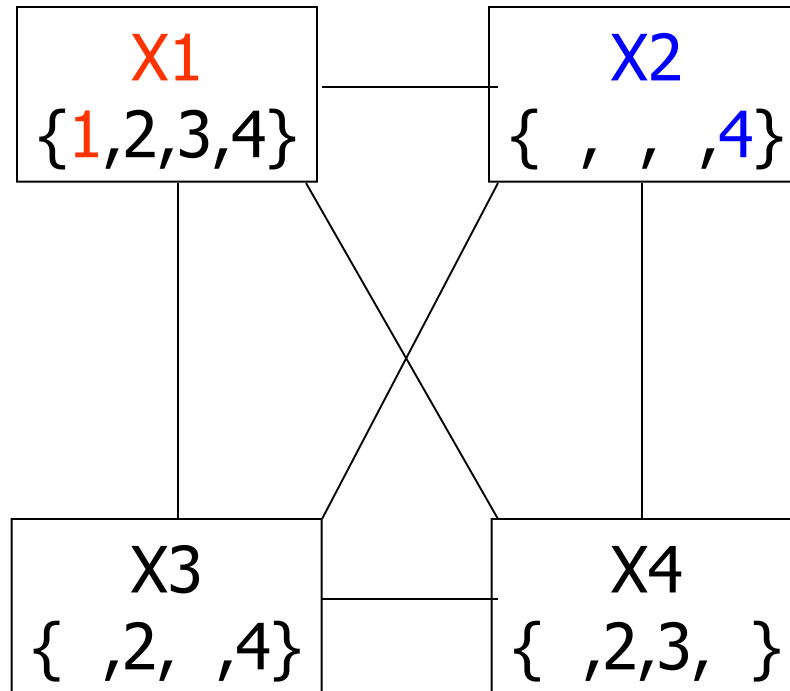
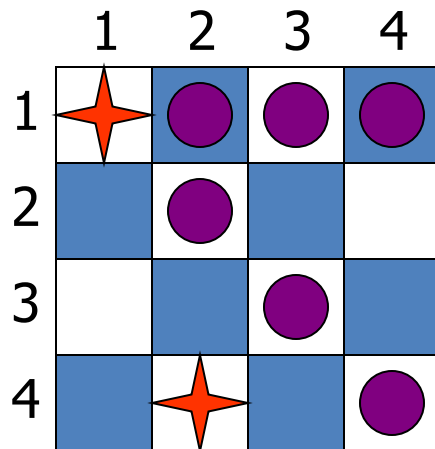
Example: 4-Queens Problem

	1	2	3	4
1				
2				
3				
4				





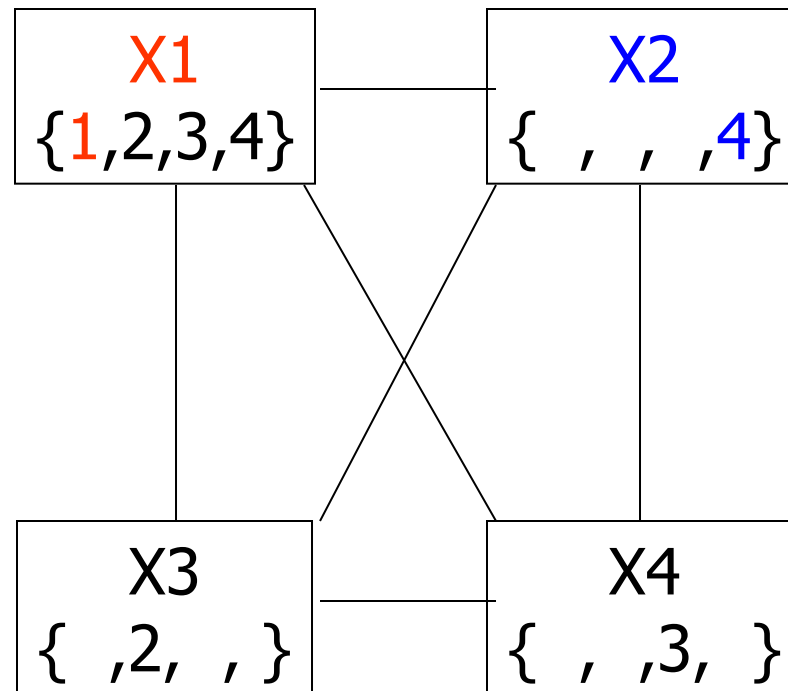
Example: 4-Queens Problem





Example: 4-Queens Problem

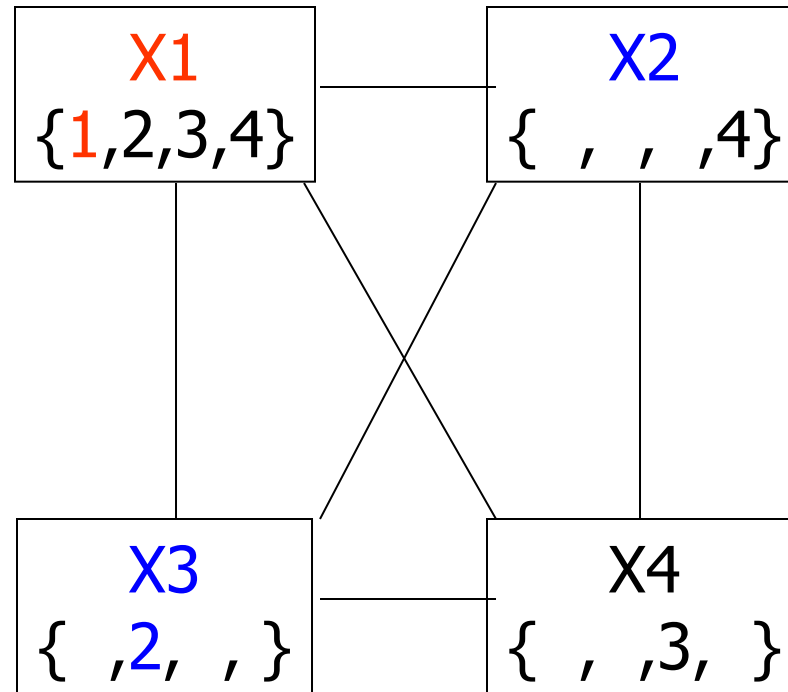
	1	2	3	4
1	★	●	●	●
2		●		●
3			●	
4		★	●	●





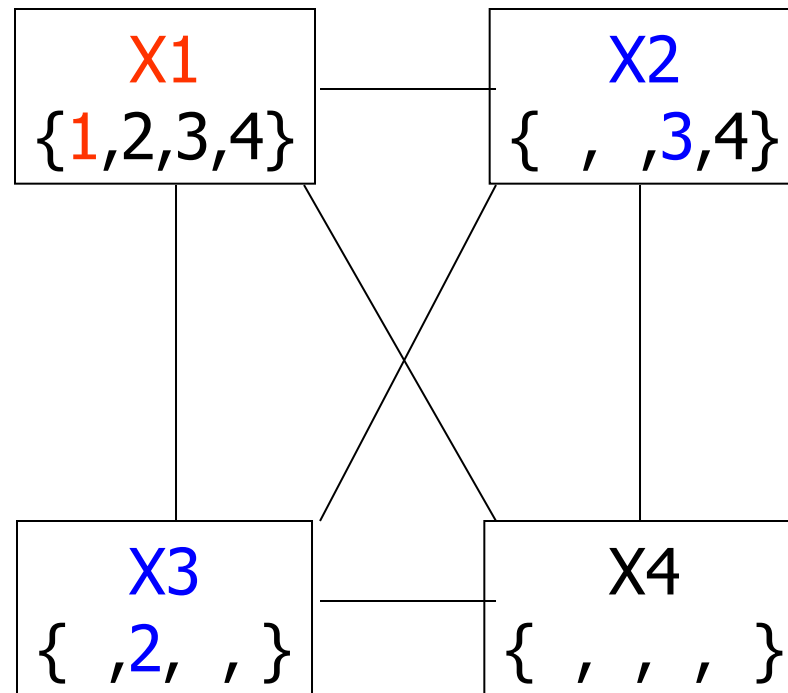
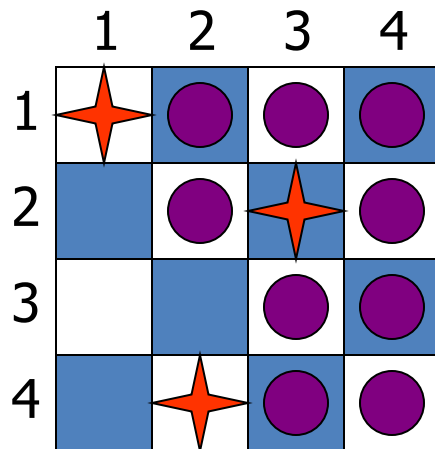
Example: 4-Queens Problem

	1	2	3	4
1	★	●	●	●
2		●	★	●
3			●	
4		★	●	●





Example: 4-Queens Problem



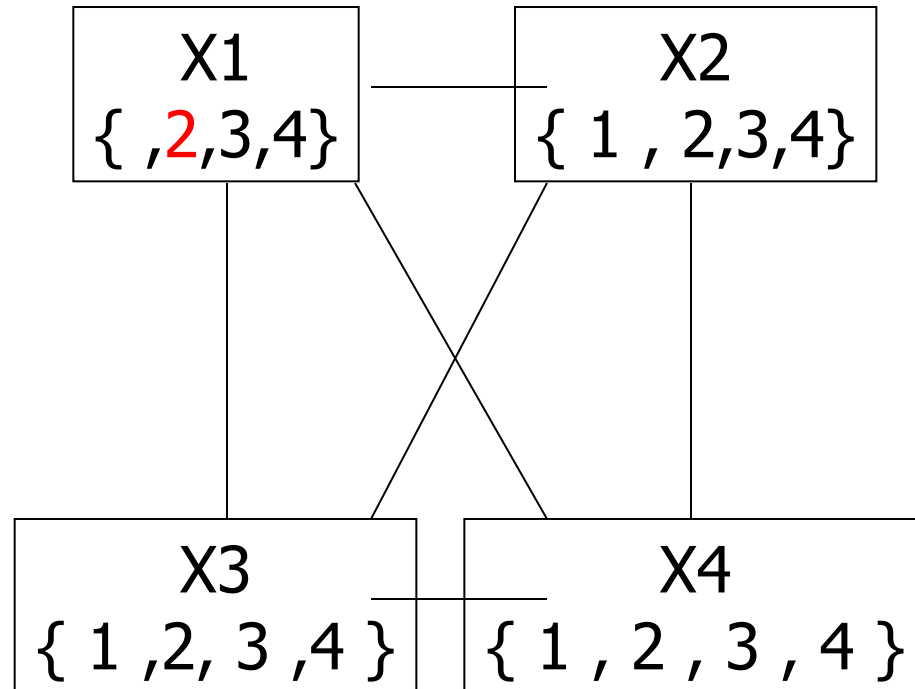
Dead End → Backtrack





Example: 4-Queens Problem

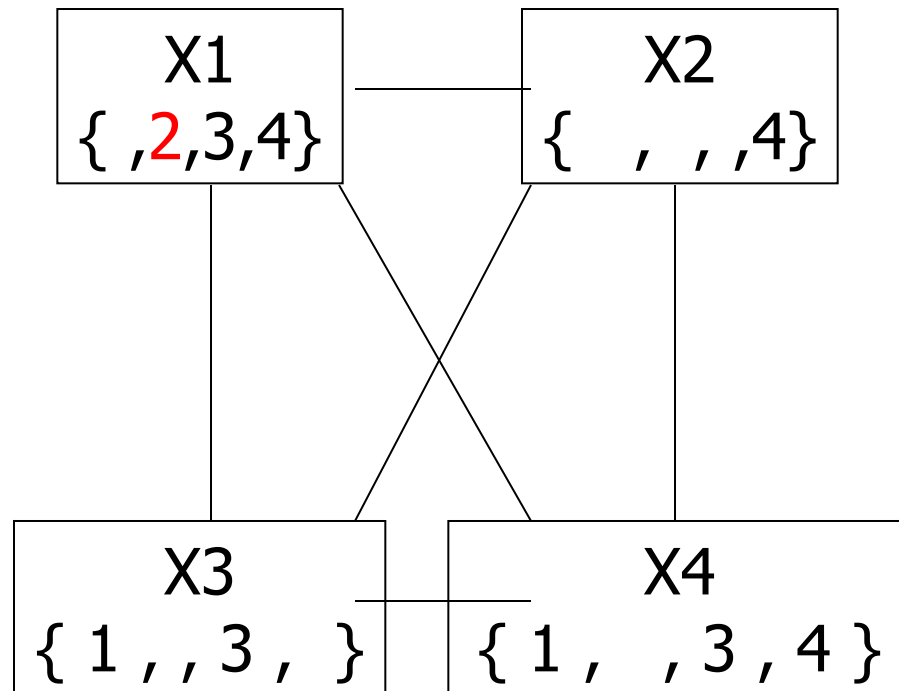
	1	2	3	4
1		●		
2	★	●	●	●
3		●		
4			●	





Example: 4-Queens Problem

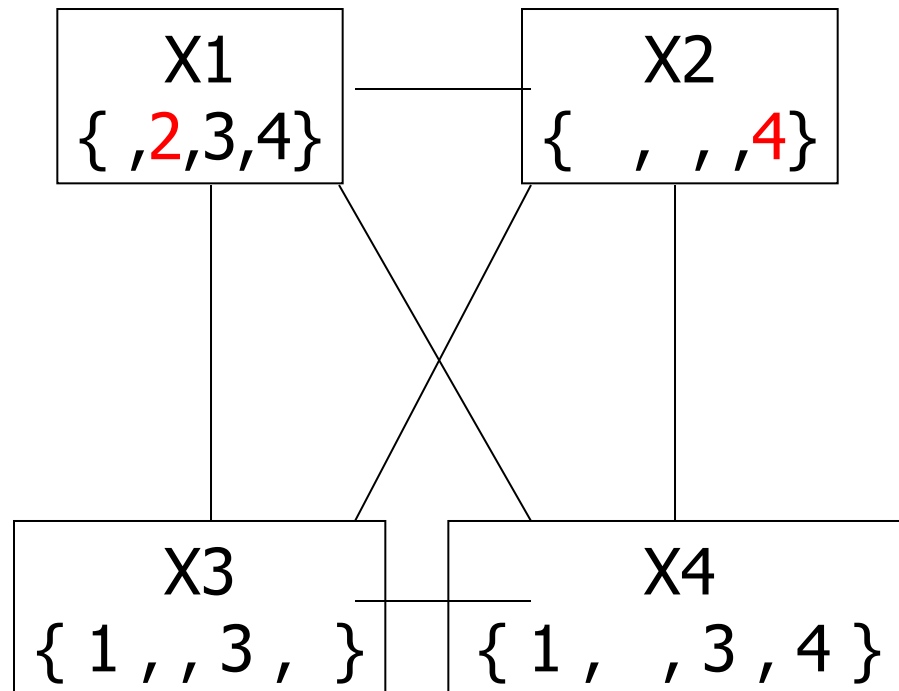
	1	2	3	4
1		●		
2	★	●	●	●
3		●		
4			●	





Example: 4-Queens Problem

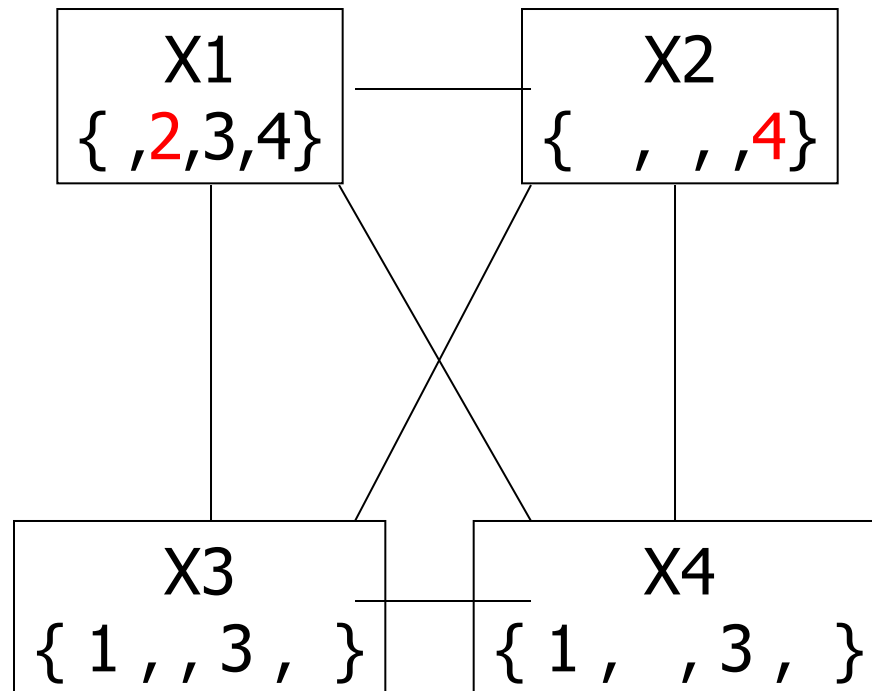
	1	2	3	4
1		●		
2	★	●	●	●
3		●	●	
4		★	●	●





Example: 4-Queens Problem

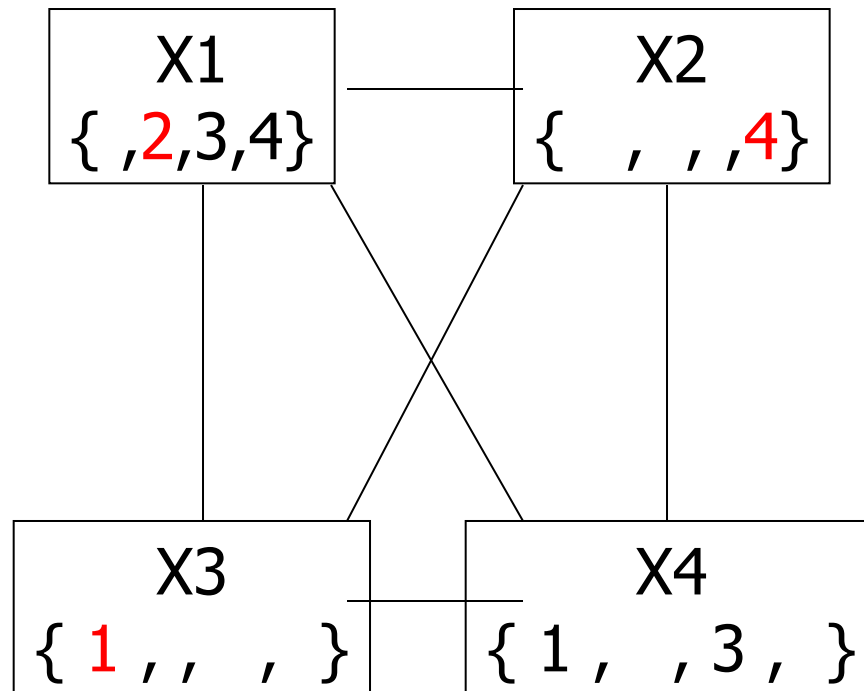
	1	2	3	4
1		●		
2	★	●	●	●
3		●	●	
4		★	●	●





Example: 4-Queens Problem

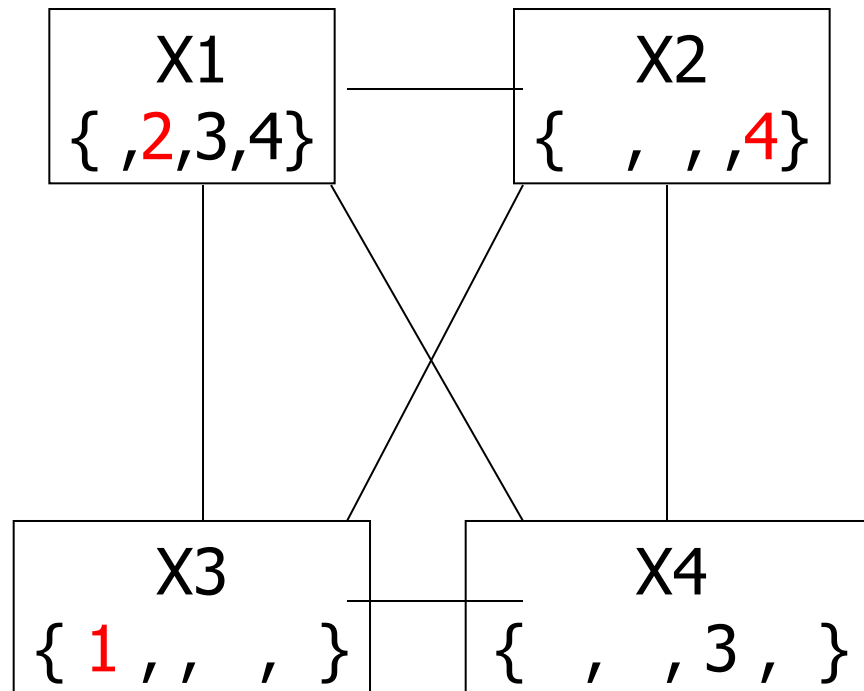
	1	2	3	4
1		●	★	●
2	★	●	●	●
3		●	●	
4		★	●	●





Example: 4-Queens Problem

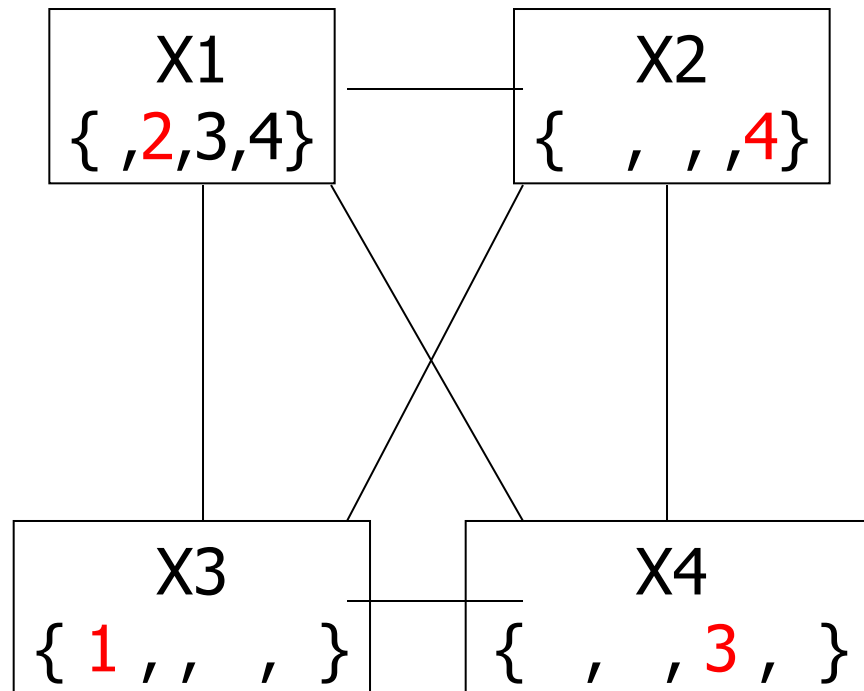
	1	2	3	4
1		●	★	●
2	★	●	●	●
3		●	●	
4		★	●	●





Example: 4-Queens Problem

	1	2	3	4
1		●	★	●
2	★	●	●	●
3		●	●	★
4	■	★	●	●



Solution !!!!





Consistency Techniques

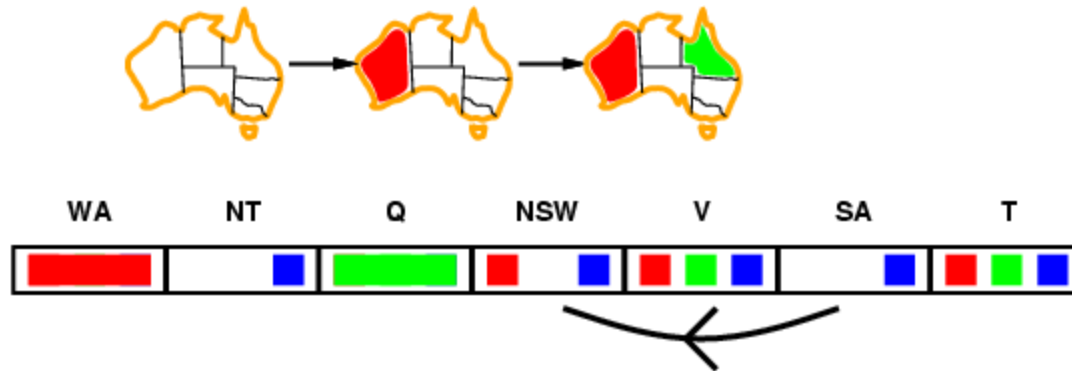
- removing inconsistent values from variables' domains
- graph representation of the CSP
- **arc consistency (AC)**





Arc consistency

- Simplest form of propagation makes each arc **consistent**
- $X \rightarrow Y$ is consistent iff
for **every** value x of X there is **some** allowed y

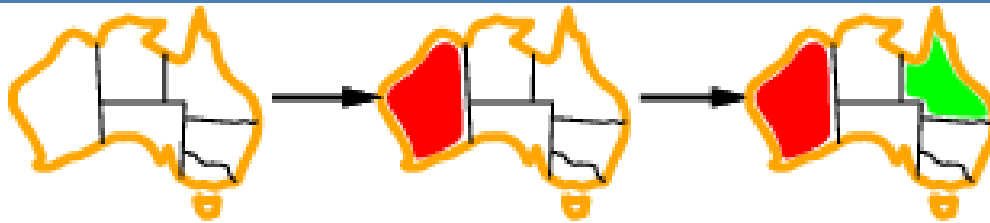


Arc from SA to NSW is consistent because SA (blue) allows NSW (red)





Arc consistency



WA

NT

Q

NSW

V

SA

T



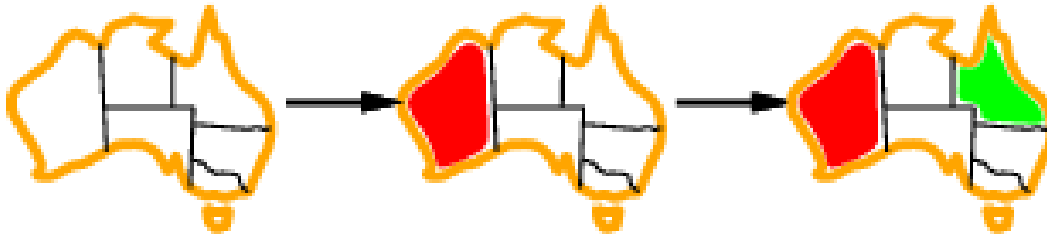
Arc from NSW to SA is **not consistent** because NSW (blue) does not allow SA (blue)

Can be made consistent by removing blue from NSW domain





Arc consistency



WA

NT

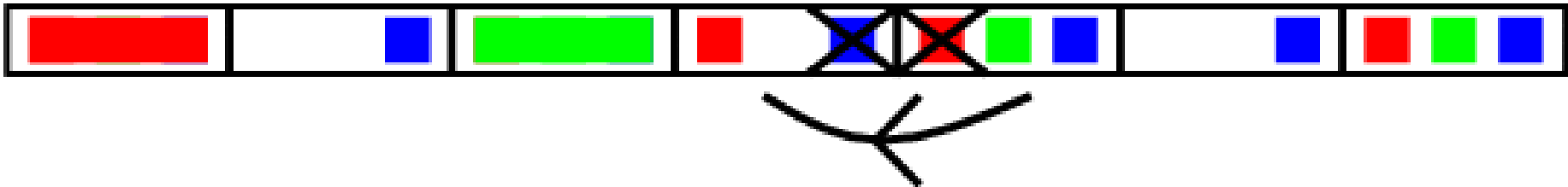
Q

NSW

V

SA

T



If X loses a value, neighbors of X need to be rechecked.

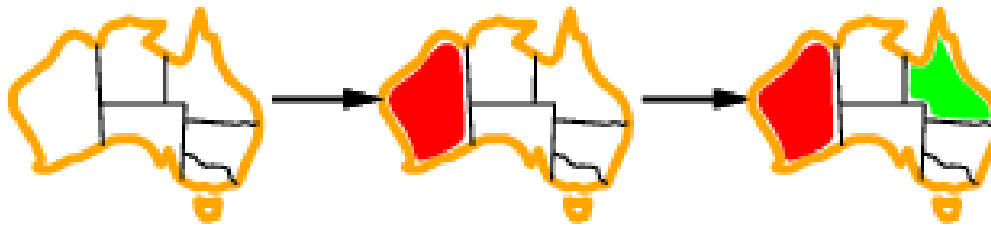


After removing blue from NSW, must check neighbors for consistency leading to removal of red from V.





Arc consistency



WA

NT

Q

NSW

V

SA

T



- If X loses a value, neighbors of X need to be rechecked
- Arc consistency detects failure earlier than forward checking
- Can be run as a preprocessor or after each assignment

After removing blue, SA domain is $\{\}$, pruning SA completely for the two assignments $\{\text{NSW}=\text{red}, \text{Q}=\text{green}\}$ and forcing backtracking.





- Is the arc (WA,NSW) consistent?
- Is the arc (NSW,WA) consistent?
- If not, what should be the domain of NSW and WA?
- Is the arc (SA,T) consistent?
- is the arc (T, SA) consistent?





Thank you



**End of
Chapter 6**

